

Lawrence Livermore National Laboratory

FPGA based Network Traffic Analysis using Traffic Dispersion Graphs

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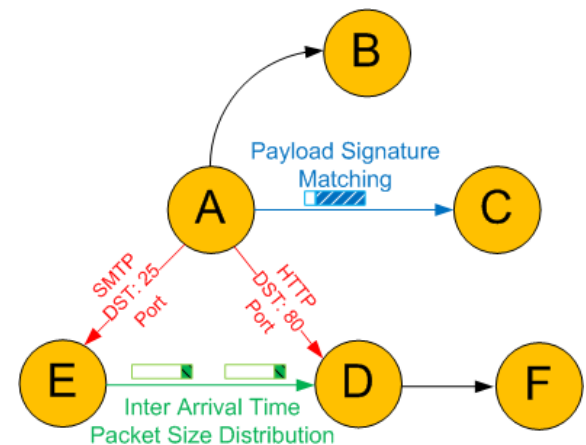
Outline

- Introduction and Motivation
- FPGA Architecture & Algorithm
- Experimental Results
- Conclusions and Future Work



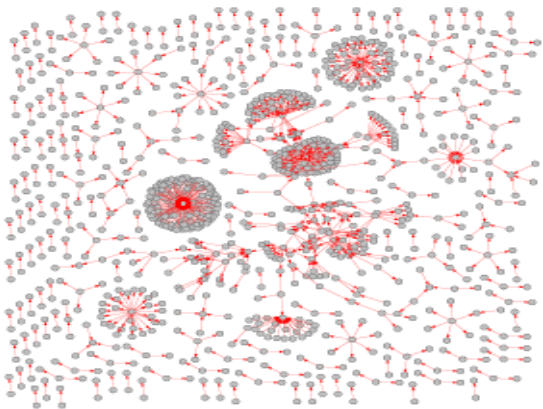
Network Traffic Analysis

- Network traffic analysis & classification is keystone in wide range of applications such as detection of network anomalies and attacks, identification of new applications, and traffic engineering.
- Traditionally it has been done using
 - **TCP destination port**
 - **Payload signature matching**
 - **Behavioral characterization of a host and its flows**

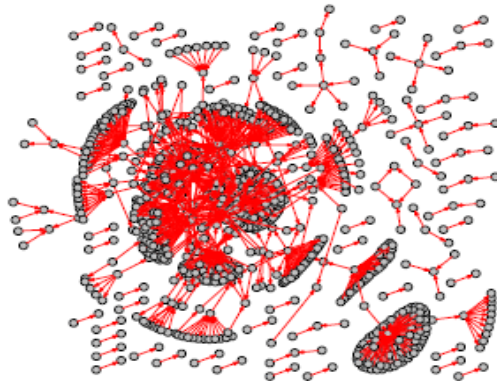


Traffic Dispersion Patterns

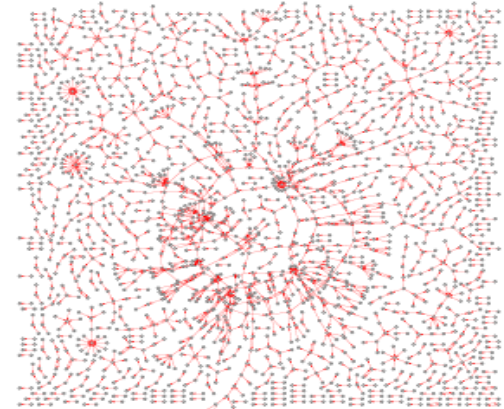
- Traffic Dispersion Patterns/Graphs (TDPs/TDGs) are a network wide extension of behavioral techniques [\[Liofotou-07\]](#).
- Communication patterns between the hosts exhibit insights into different kinds of applications and anomalies in a network.



HTTP



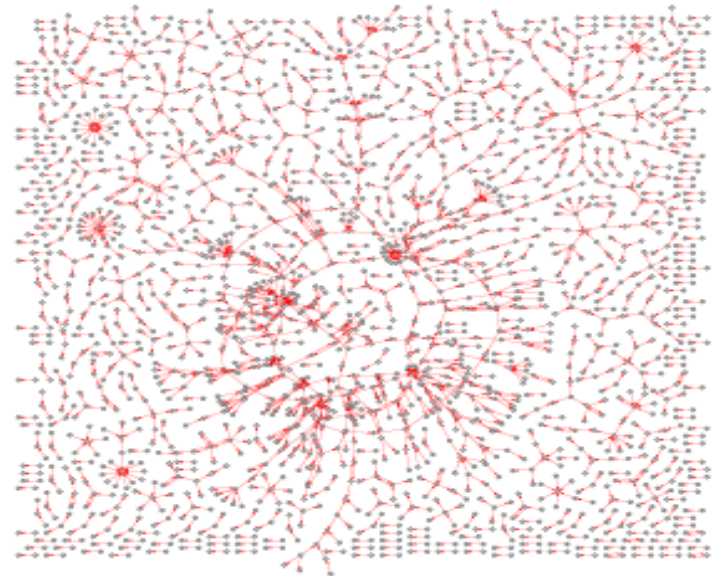
SMTP



WinMX (Peer-to-Peer)

Applications

- Isolating anomalous hosts and sub-networks
- Detecting the spread of worm and viruses.
- Characterizing unknown activity/application.
- Flagging suspicious activity without examining the content.



WinMX (Peer-to-Peer)

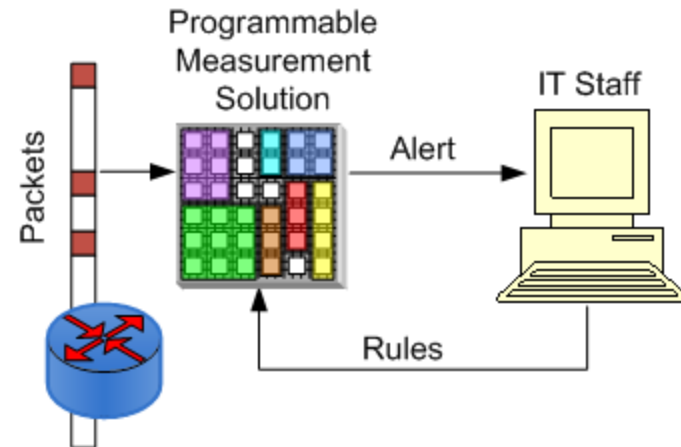
Motivation

- TDPs are traditionally processed offline (non-streaming)
 - Full view of the graph: Complete flexibility
 - Slow and resource intensive
- Real-time threats get missed out.
- The question of processing TDPs in real-time has mostly been overlooked.



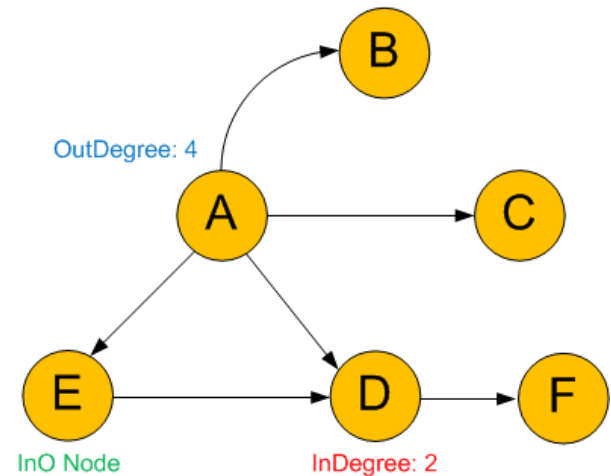
Goals

- An autonomous passive system that can detect anomalous behaviors in streaming traffic in real-time.
- Isolates or alarms the IT staff about the threats.
- Programmable to varying network monitoring needs.
- Scalable to millions of edges and new traffic features.
- Has a defined accuracy budget.



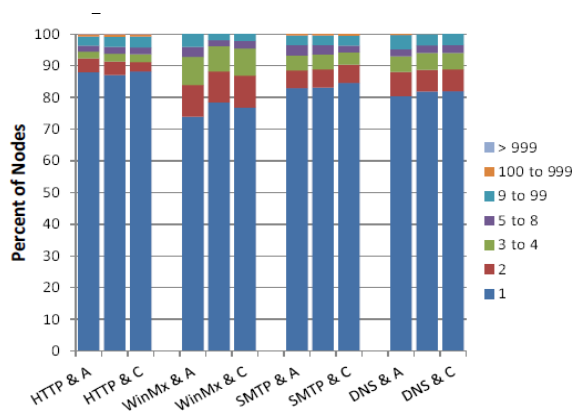
Quantizing Traffic Dispersion Patterns

- Dispersion patterns can be classified as
 - Offline: requires knowledge of complete graph. For e.g. depth of communication
 - Online: can be computed in streaming traffic. For e.g. connectivity of an IP address
- InDegree/OutDegree: The number of incoming/outgoing connections made to/from an IP address/Node.
- InO: The number of hosts that have both incoming and outgoing connections.

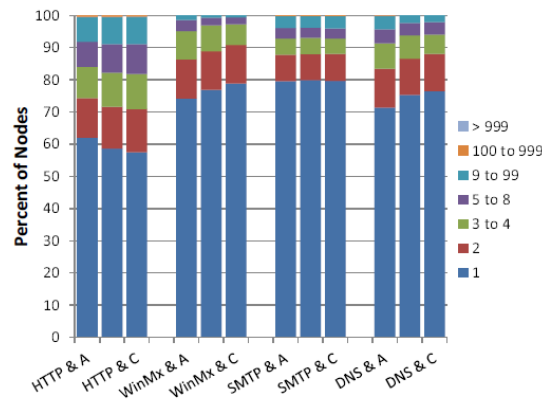


Defining Rules

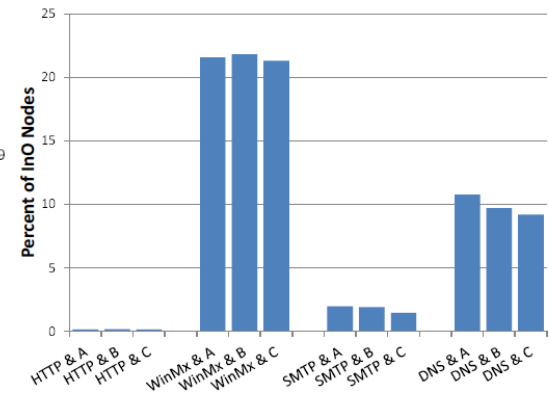
- Simulated backbone CAIDA traces having known P2P activity.



InDegree Distributions



OutDegree Distributions



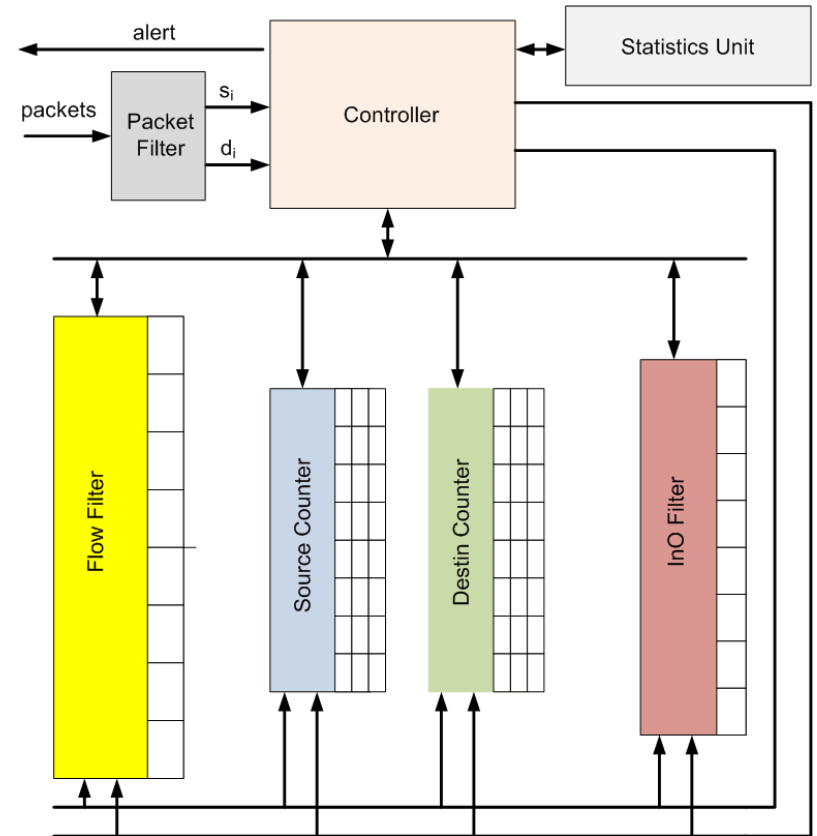
InO Distributions

- P2P Detection Rule: A high InO with a low InDegree and OutDegree
- Port-Scan Rule: High InDegree and/or High OutDegree



The Measurement Solution

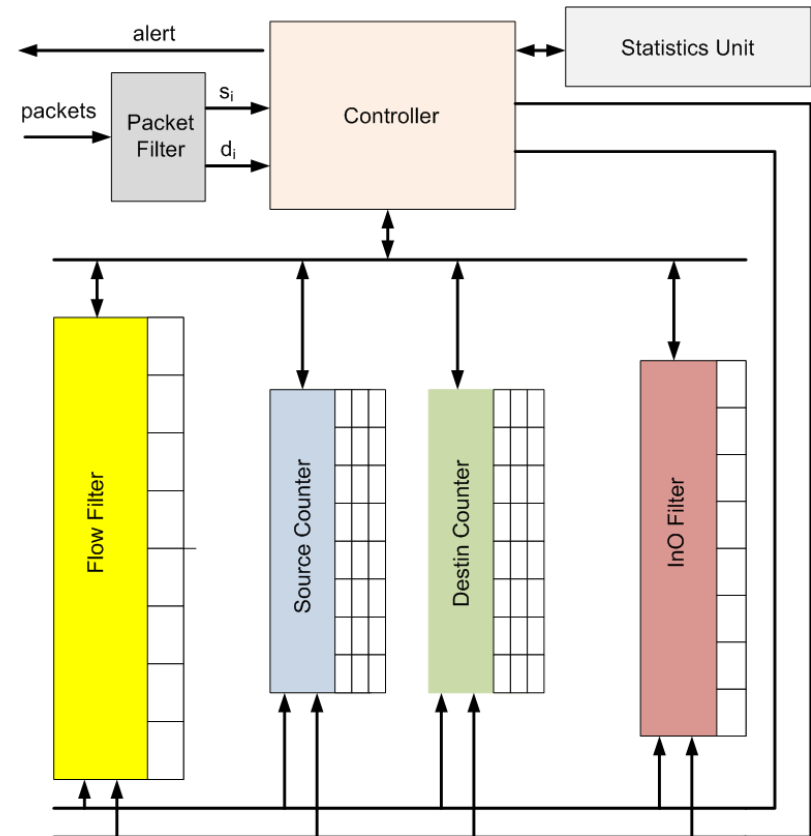
- Mapped on FPGAs that are an ideal platform for the high programmability and throughput requirements of the application.
- Configurable packet filter.
- Bloom Filters to check presence of an incoming IP address.
- Bloom Counters to keep the InDegree and OutDegrees.
- Count values and programmable rules in Statistics Unit.



FPGA Mapped Architecture

The Measurement Algorithm

- Incoming packet checked for uniqueness in the Flow-filter.
- The source-IP and destination-IP are used to update the degree-counts.
- Statistics unit is updated.
- The IPs are reverse checked for deducing InO behavior.
- A found InO IP is checked in InO-Filter and InO count updated

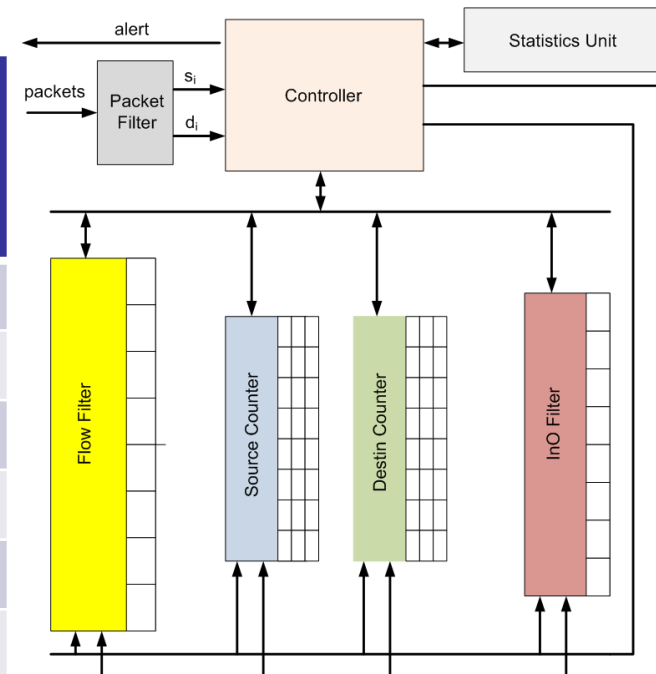


FPGA Mapped Architecture

Design Space Exploration

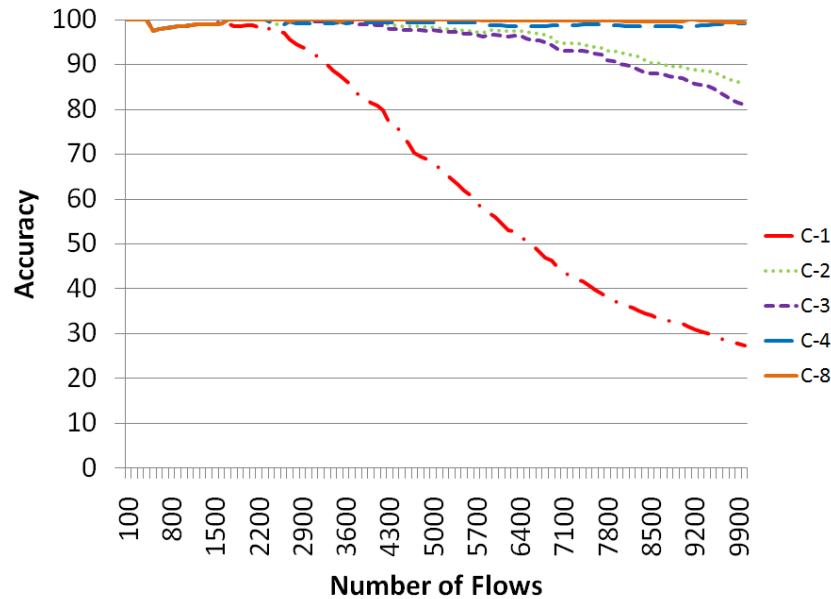
Config	Flow x8	InO x8	Degree x2x8x8	Total (Kb)	Node Count Accuracy	InO Count Accuracy
C-1	16	16	2	512	71.78	27.35
C-2	8	8	4	640	88.31	85.75
C-3	16	8	4	704	92.14	81.07
C-4	8	8	8	1152	90.13	99.18
C-8	32	1	8	1312	96.95	99.34
C-9	32	16	16	2432	99.98	99.67

**Architectural Exploration
(values in Kb)**

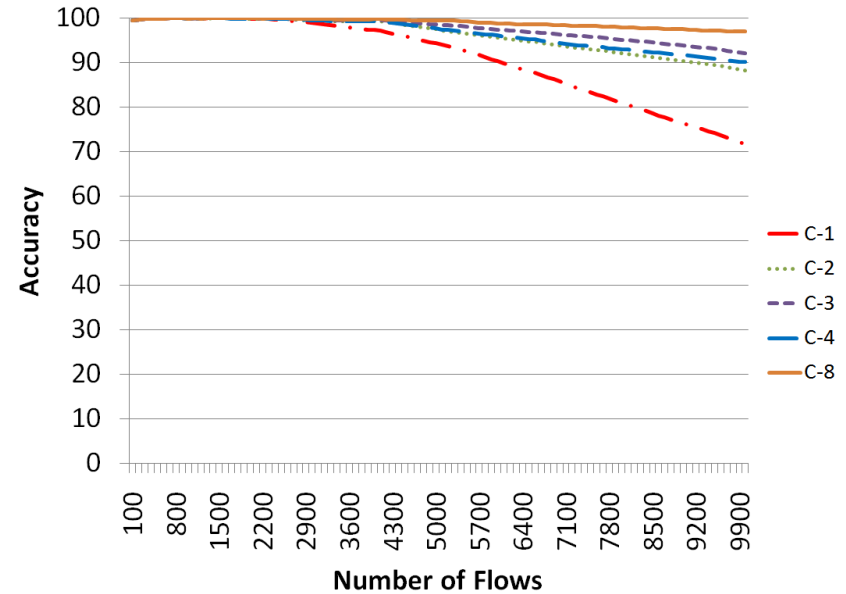


FPGA Mapped Architecture

Design Accuracy



Node Count Accuracy



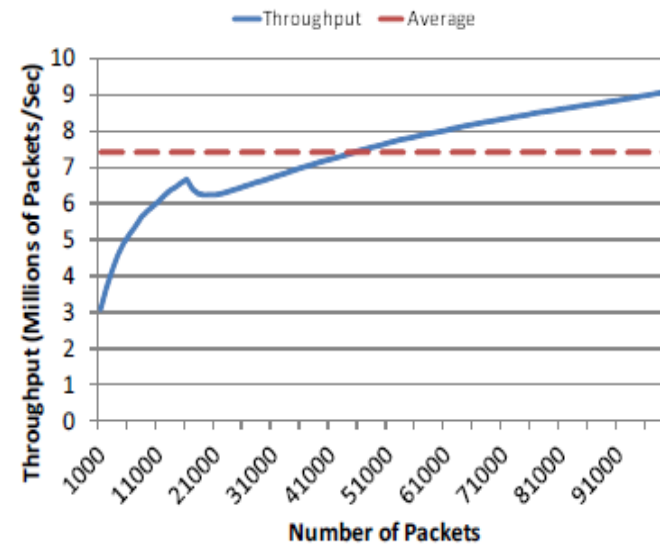
In0 Count Accuracy

- The work utilized datasets from CAIDA, WIDE and ISI/USC for empirical evaluations.
- The work discusses analytical and empirical accuracy studies.
- Results show very high amounts of accuracies for up to 10k flows.
- Suitable for detecting threats like Port-Scan.

Implementation Results & System Throughput

- Prototyped C-8 on Virtex-5, XC5VLX50t device.
- Xilinx ISE 11.1
- Ample room in the device for future logic enhancements.
- Average throughput of 7.5 million packets/second.

Metric	Value	Device Utilization
IO Pins	69	14.4%
Number of occupied slices	754	10%
Number of Block-RAMs	44	73%
Clock (MHz)	107	-



Port-Scan Detection Accuracy

- Very high detection accuracy of both scanners and victim identification.

		Observed	
		P	N
Actual	Scanners		
	P'	4	0
N'	0	33972	

Scanner Identification

F1 Score: 1.0

		Observed	
		P	N
Actual	Victims		
	P'	22	0
N'	2	2303	

Scanned Victim Identification

F1 Score: 0.932

- Used USC-LANDER **IUSCISI-06I** and CAIDA traces **ICAIDA-03I**.

Peer-to-Peer Detection Accuracy

- Programmed Rule: 10% of InO Nodes and Low InDegree and OutDegrees.
- Port based filtering.
- Alarm raised in less than 1400 flows at the correct port (port-6677)
- 1400 flows represent a very high system accuracy design point.
- Useful for classifying P2P activity on a new port.



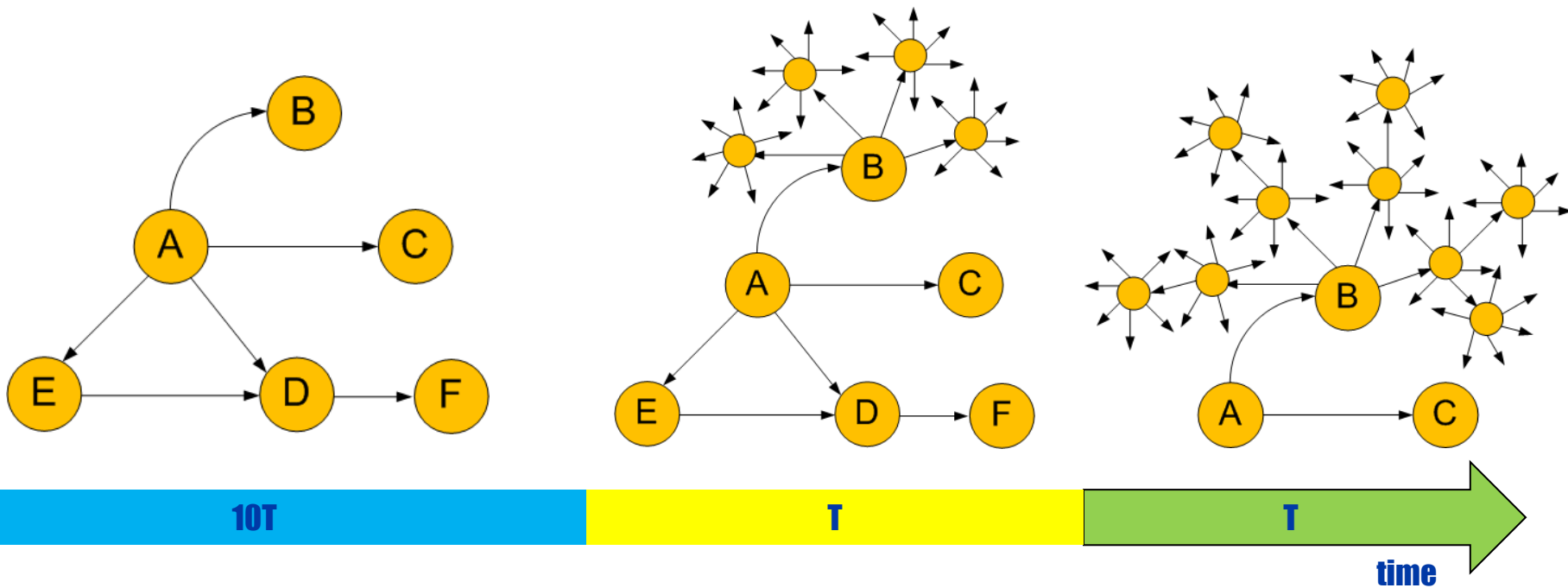
Conclusion

- A real-time FPGA based solution for processing streaming TDPs.
- Configurable with high processing speeds.
- Real-time detection of Port-Scan and Peer-to-Peer activities.
- Reducing noise in capturing interesting traffic.
- Easily extensible to include temporal rules and other characterization measures.
- Open Source under GNU General Public License:
https://computation.llnl.gov/casc/dcca-pub/dcca/Data-centric_architecture.html



Future Directions

- Including temporal relations in rule sets
- Isolating sudden changes from gradual shifts
- For e.g. A new Worm, Spammer, etc



References

- **[Iliofotou-07]** M. Iliofotou, P. Pappu, M. Faloutsos, M. Mitzenmacher, S. Singh, and G. Varghese, “Network monitoring using traffic dispersion graphs (TDGS),” in *Proceedings of the 7th ACM SIGCOMM Internet Measurement Conference, 2007*, pp. 315–320.
- **[USCISI-06]** Internet Addresses Survey dataset, PREDICT ID: USC-LANDER/attack-tcpsyn-20061106. Traces taken 2006-11-06. Provided by the USC/LANDER project <<http://www.isi.edu/ant/lander>>.
- **[CAIDA-03]** CAIDA OC48 Trace Project 20030115-100000-0-anon.pcap
[https://imdc.datcat.org/data/1-3N3T-4=20030115-100000-0-anon.pcap+\(60+min\)](https://imdc.datcat.org/data/1-3N3T-4=20030115-100000-0-anon.pcap+(60+min))